
Beyond Single-Deletion Correcting Codes: Substitutions and Transpositions

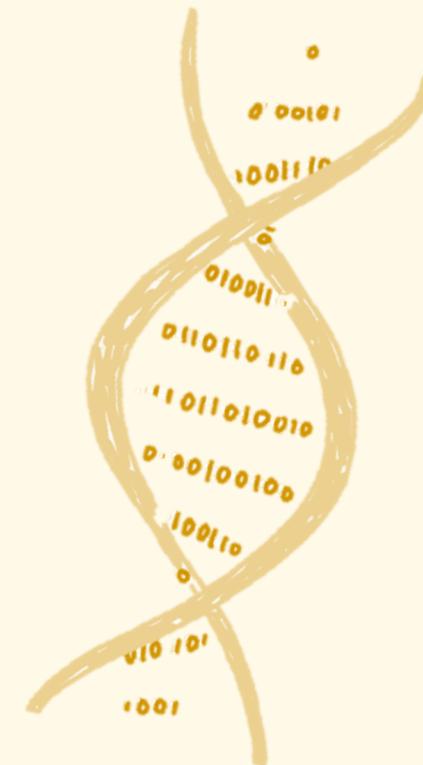
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Three types of common errors

Substitution

Deletion(insertion)

Transposition



Substitution

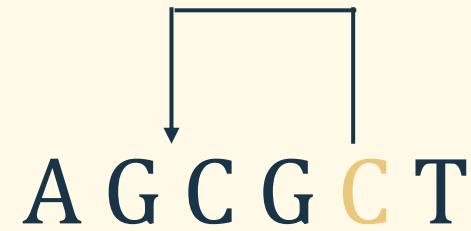
A G C G **C** T
↓
A G C G **T** T

Deletion(insertion)

A G C G C T

A G C G T

Transposition

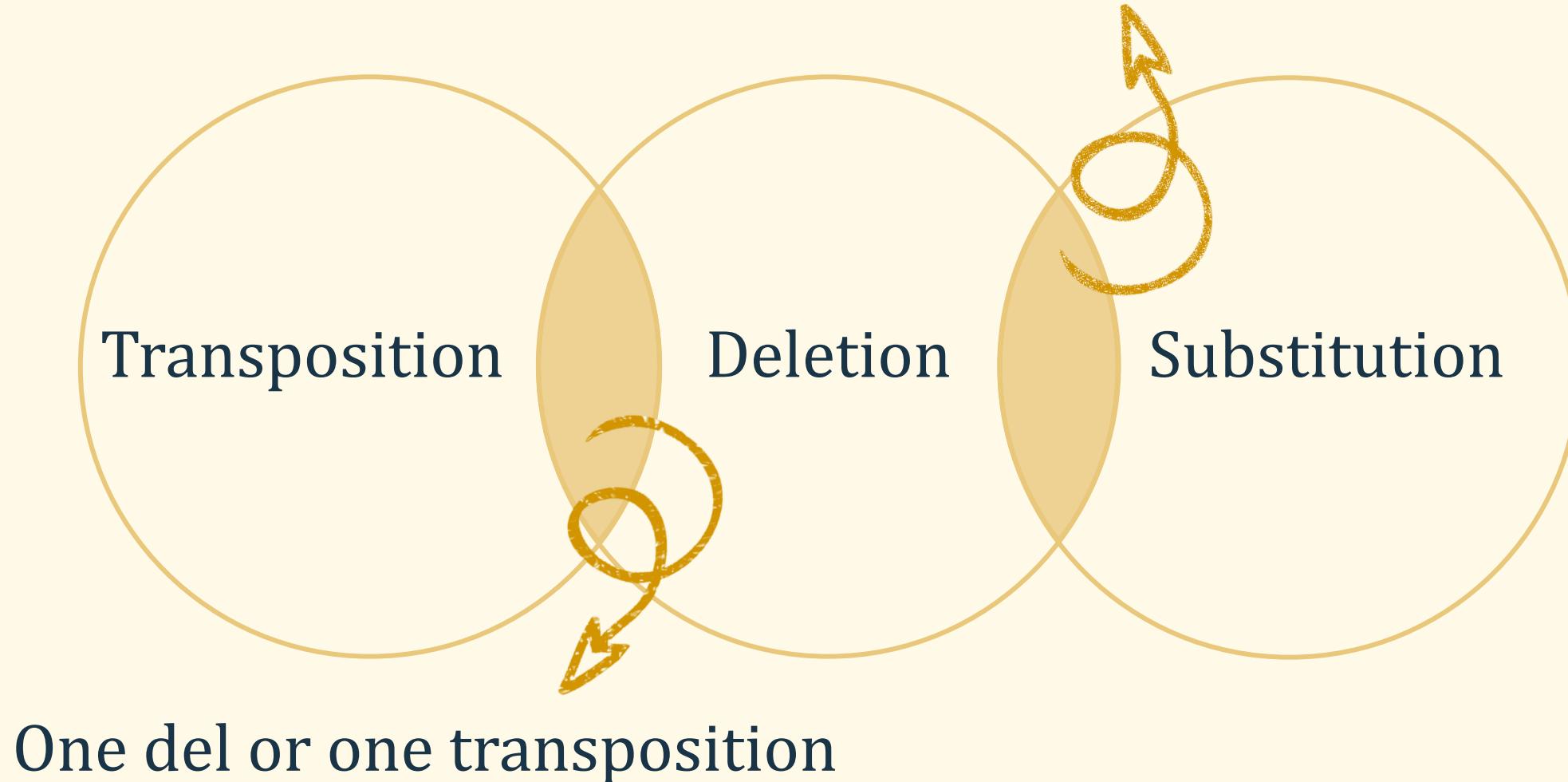


A G C G C T



A G C C G T

How about the interplay?



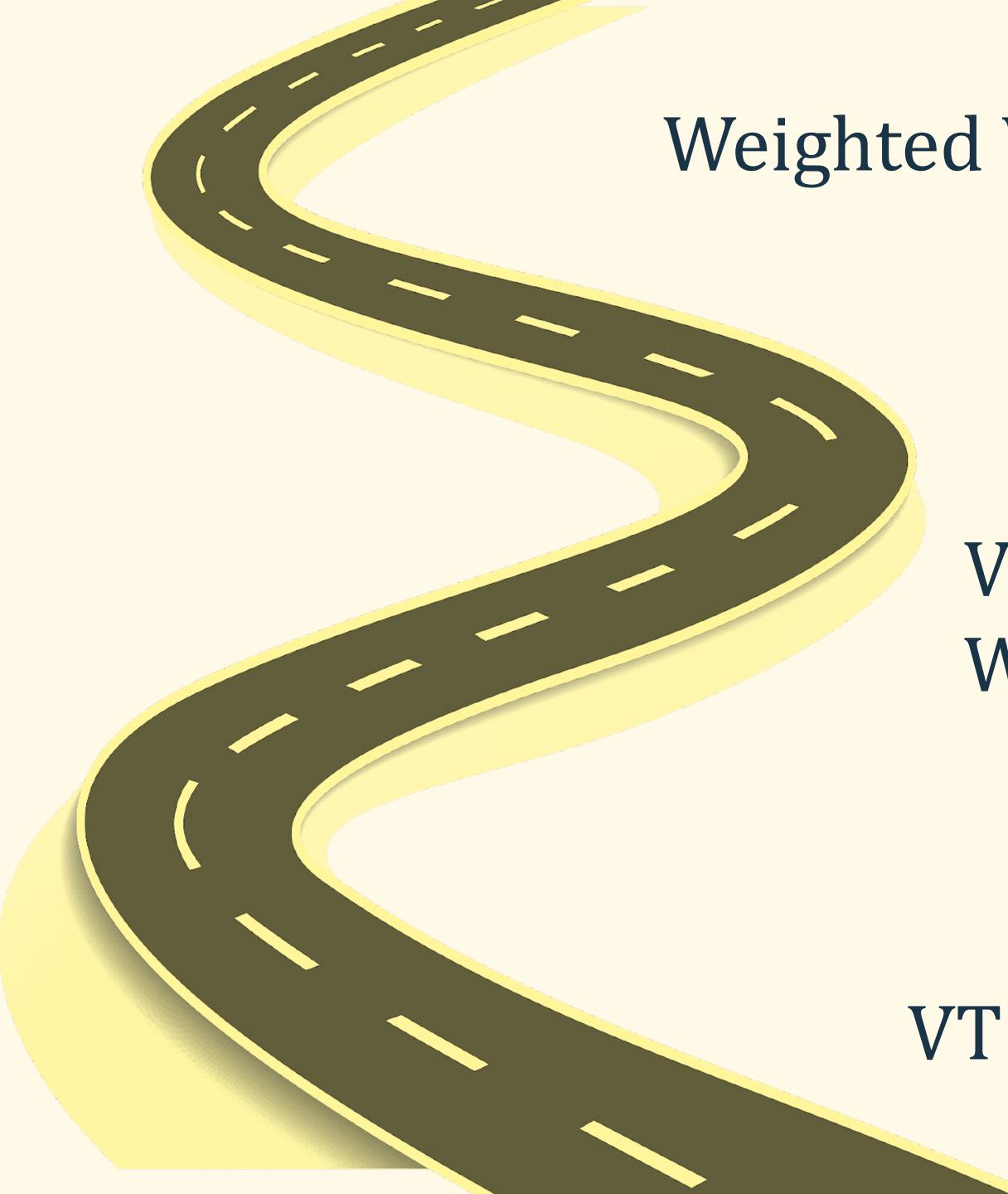
one del and one sub
one del or one sub

Our Results

Alphabet	Error type	Redundancy
q	One del or one sub (edit error)	$\log n + O_q(\log \log n)$
2	One del or one adjacent trans	$\log n + O(\log \log n)$
1	One del AND one sub	$4\log n + O(\log \log n)$ List of size 2

Our Results

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q	One del or one sub (edit error)	$\log n + O_q(\log \log n)$



Weighted VT sketch

VT sketch: ternary
Why it fails

VT sketch: binary



Weighted VT sketch

VT sketch: ternary
Why it fails

VT sketch: binary

Correcting one edit error: binary

VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

If substitution: $x = 00101$ $f(x) = 8$

VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$f(x) = 8$$

$$y = 00111 \qquad f(y) = 12$$

$$f(y) - f(x) = \sum_1^n i \cdot (y_i - x_i) = e(y_i - x_i) = 4$$

Position of substitution



VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$x = 001\textcolor{brown}{0}1 \quad f(x) = 8$$

$$y = 001\textcolor{brown}{1}1 \quad f(y) = 12$$

Correcting one edit error: deletion

VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$f(x) = 8$$

$$y = 0001 \qquad \qquad f(y) = 4$$

$$f(x) - f(y) = \sum_d^{n-1} y_i + dx_d = 4$$

VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$f(x) = 8$$

$$y = ? \textcolor{brown}{0} \textcolor{brown}{0} \textcolor{brown}{0} \textcolor{brown}{1} \qquad f(y) = 4$$

$$f(x) - f(y) = \sum_d^{n-1} y_i + dx_d = 4$$



VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$x = 001\textcolor{brown}{0}1 \quad f(x) = 8$$

$$y = 000?\textcolor{brown}{1} \quad f(y) = 4$$

$$f(x) - f(y) = \sum_d^{n-1} y_i + dx_d = 4$$

VT sketch

$$\text{VT sketch: } f(x) = \sum_{i=1}^n i \cdot x_i$$

$\log n + 2$ bits!



Weighted VT sketch

VT sketch: ternary
Why it fails

VT sketch: binary

VT code for ternary?

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$x = 0 2 1 0 1 2$$

$$f(x) = 24$$

VT code for ternary: substitution

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$f(x) = 24$$

$$y = 011012$$

$$f(y) = 22$$

$$f(y) - f(x) = e(y_i - x_i) = 2$$

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$x = 2 \ 1 \ 1 \ 0 \ 1 \ 2$$

$$x = 0 \ \textcolor{brown}{2} \ 1 \ 0 \ 1 \ 2$$

$$f(x) = 24$$

$$y = \ 0 \ 1 \ 1 \ 0 \ 1 \ 2$$

$$f(y) = 22$$

$$f(y) - f(x) = e(y_i - x_i) = 2$$

$$\#0 = 2$$

$$\#1 = 2$$

$$\#2 = 2$$



$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$x = 0 \ 2 \ 1 \ 0 \ 1 \ 2 \qquad f(x) = 24$$

$$y = 0 \ 1 \ 1 \ 0 \ 1 \ 2 \qquad f(y) = 22$$

VT code for ternary: deletion of 0

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$f(x) = 24$$

$$y = 0 2 1 1 2$$

$$f(y) = 21$$

$$f(x) - f(y) = \sum_d^{n-1} y_i = 3$$



$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$x = 0 \ 2 \ 1 \ \textcolor{brown}{0} \ 1 \ 2 \qquad \qquad f(x) = 24$$

$$y = 0 \ 2 \ 1 \ ? \ \textcolor{brown}{1} \ \textcolor{brown}{2} \qquad \qquad f(y) = 21$$

$$f(x) - f(y) = \sum_d^{n-1} y_i = 3$$

Position of deletion

VT code for ternary: deletion of 1

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$f(x) = 24$$

$$y = 0 2 0 1 2$$

$$f(y) = 18$$

$$f(x) - f(y) = \sum_d^{n-1} y_i + d = 6$$

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$x = 1 \ 0 \ 2 \ 0 \ 1 \ 2$$

$$x = 0 \ 2 \ \textcolor{brown}{1} \ 0 \ 1 \ 2$$

$$f(x) = 24$$

$$y = \ 0 \ 2 \ 0 \ 1 \ 2$$

$$f(y) = 18$$

$$f(x) - f(y) = \sum_d^{n-1} y_i + d = 6$$

$x = 1 \textcolor{brown}{0} \textcolor{brown}{2} 0 1 2$

$x = 0 \textcolor{brown}{2} \textcolor{brown}{1} 0 1 2$

$$\sum_d^{n-1} y_i + d$$

$$\sum_d^{n-1} y_i + d$$

$x = 1 \boxed{0} \boxed{2} 0 1 2$ $x = 0 \boxed{2} \boxed{1} 0 1 2$

$$\sum_d^{n-1} y_i + \textcolor{brown}{d}$$

$$\sum_d^{n-1} y_i + \textcolor{brown}{d}$$

Fails when there is a chunk with an average 1



Weighted VT sketch

VT sketch: ternary
Why it fails

VT sketch: binary



Bias the weight!

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$f(x) = \sum_{i=1}^n i \cdot w(x_i)$$

$$w(0) = 0$$

$$w(1) = 1$$

$$w(2) = 2 \log n$$

$$x = 1 \textcolor{brown}{0} \textcolor{brown}{2} \dots$$

$$f(x) = \sum_{i=1}^n i \cdot w(x_i)$$

$$w(0) = 0$$

$$w(1) = 1$$

$$\textcolor{brown}{w}(2) = 2 \log n$$

$$f(x) = 1 + 6 \log n$$

$$x = \textcolor{brown}{0} \textcolor{brown}{2} 1 \dots$$

$$f(x) = 4 \log n + 3$$

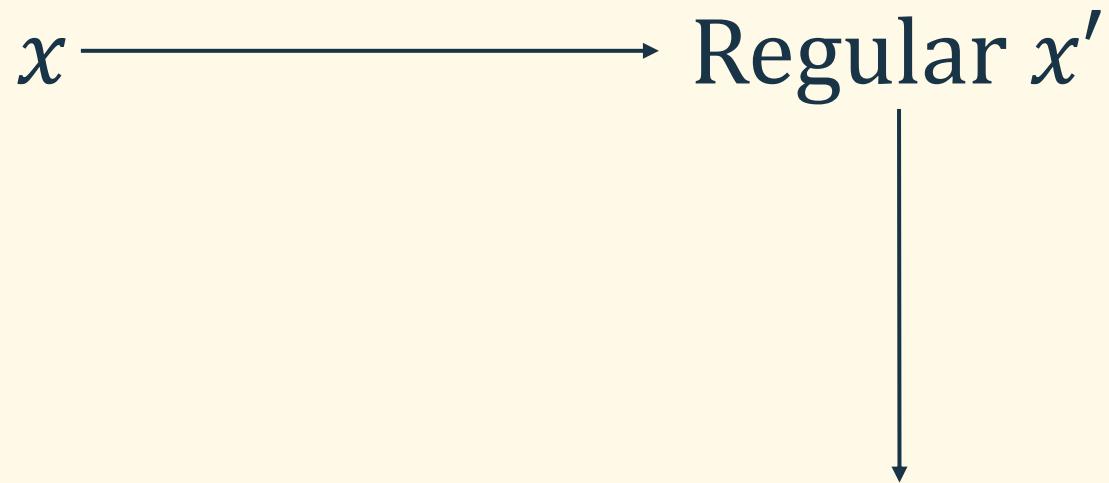
As long as the chunk of avg 1 has length $< 2 \log n$

$$x = 1 \underbrace{0 \ 2 \ 0 \ 2 \ 0 \ 2 \ 0 \ 2 \ 0 \ 2 \ 2 \dots}_{< 2 \log n}$$

Run of 0 after deleting all 2's $< \log n$

$$x = 1 \underbrace{0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 2 \dots}_{< \log n}$$

Weighted VT: encoding



Weighted VT of x'

$\log n + \log \log n$ redundancy!

More in paper

1. Extends to general q
2. Binary code correcting one del and one sub
3. Binary code correcting one del and one adjacent trans