
Beyond Single-Deletion Correcting Codes:

Substitutions and Transpositions

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Three types of common errors

Substitution

Deletion(insertion)

Transposition



Substitution

A G C G C T



A G C G T T

Deletion(insertion)

A G C G C T

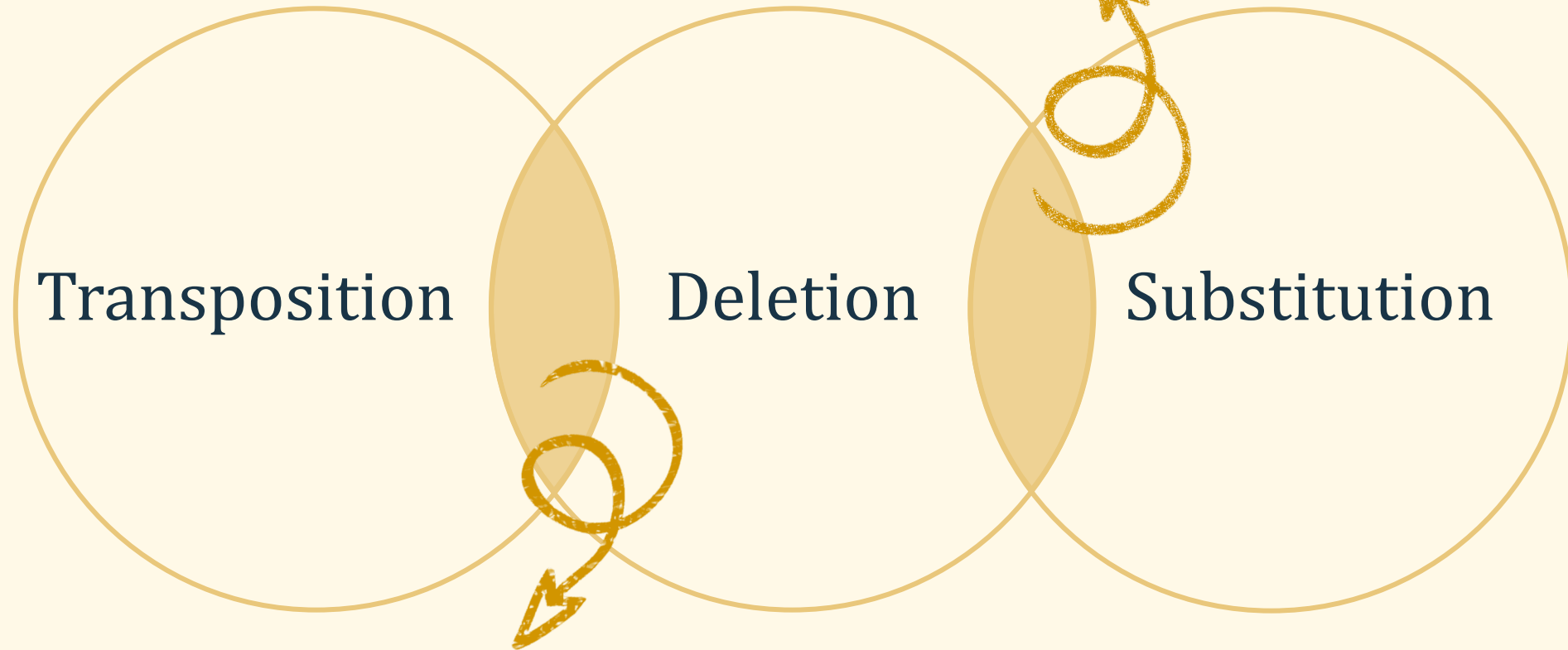
A G C G T

A G C G C T

A G C C G T

Transposition

How about the interplay?



one del and one sub
one del or one sub

One del or one transposition

Our Results

Alphabet	Error type	Redundancy
q	One del or one sub (edit error)	$\log n + O_q(\log \log n)$
2	One del or one adjacent trans	$\log n + O(\log \log n)$
1	One del AND one sub	$4\log n + O(\log \log n)$ List of size 2

Our Results

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Weighted VT sketch

VT sketch: ternary
Why it fails

VT sketch: binary



Weighted VT sketch

VT sketch: ternary
Why it fails

VT sketch: binary

Correcting one edit error: binary

VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

If substitution:

$$x = 0\ 0\ 1\ 0\ 1$$

$$f(x) = 8$$

VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$f(x) = 8$$

$$y = 0 \ 0 \ 1 \ 1 \ 1 \qquad f(y) = 12$$

$$f(y) - f(x) = \sum_1^n i \cdot (y_i - x_i) = e(y_i - x_i) = 4$$

Position of substitution



VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$x = 0 \ 0 \ 1 \ 0 \ 1 \qquad f(x) = 8$$

$$y = 0 \ 0 \ 1 \ 1 \ 1 \qquad f(y) = 12$$

Correcting one edit error: deletion

VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$f(x) = 8$$

$$y = 0\ 0\ 0\ 1$$

$$f(y) = 4$$

$$f(x) - f(y) = \sum_d^{n-1} y_i + dx_d = 4$$

VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$f(x) = 8$$

$$y = ? \text{ 0 0 0 1}$$

$$f(y) = 4$$

$$f(x) - f(y) = \sum_d^{n-1} y_i + dx_d = 4$$



VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$$x = 0 \ 0 \ 1 \ 0 \ 1 \qquad f(x) = 8$$

$$y = 0 \ 0 \ 0 \ ? \ 1 \qquad f(y) = 4$$

$$f(x) - f(y) = \sum_d^{n-1} y_i + dx_d = 4$$

VT sketch

VT sketch: $f(x) = \sum_{i=1}^n i \cdot x_i$

$\log n + 2$ bits!



Weighted VT sketch

VT sketch: ternary
Why it fails

VT sketch: binary

VT code for ternary?

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$x = 0 \ 2 \ 1 \ 0 \ 1 \ 2$$

$$f(x) = 24$$

VT code for ternary: substitution

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$f(x) = 24$$

$$y = 0 \ 1 \ 1 \ 0 \ 1 \ 2$$

$$f(y) = 22$$

$$f(y) - f(x) = e(y_i - x_i) = 2$$

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$\#0 = 2$$

$$\#1 = 2$$

$$\#2 = 2$$

$$x = 2 \ 1 \ 1 \ 0 \ 1 \ 2$$

$$x = 0 \ 2 \ 1 \ 0 \ 1 \ 2$$

$$f(x) = 24$$

$$y = 0 \ 1 \ 1 \ 0 \ 1 \ 2$$

$$f(y) = 22$$

$$f(y) - f(x) = e(y_i - x_i) = 2$$



$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$x = 0 \text{ } 2 \text{ } 1 \text{ } 0 \text{ } 1 \text{ } 2$$

$$f(x) = 24$$

$$y = 0 \text{ } 1 \text{ } 1 \text{ } 0 \text{ } 1 \text{ } 2$$

$$f(y) = 22$$

VT code for ternary: deletion of 0

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$f(x) = 24$$

$$y = 0 \ 2 \ 1 \ 1 \ 2$$

$$f(y) = 21$$

$$f(x) - f(y) = \sum_{d=1}^{n-1} y_d = 3$$



$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$x = 0 \ 2 \ 1 \ 0 \ 1 \ 2$$

$$f(x) = 24$$

$$y = 0 \ 2 \ 1 \ ? \ 1 \ 2$$

$$f(y) = 21$$

$$f(x) - f(y) = \sum_d^{n-1} y_d = 3$$

Position of deletion

VT code for ternary: deletion of 1

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$f(x) = 24$$

$$y = 0 \ 2 \ 0 \ 1 \ 2$$

$$f(y) = 18$$

$$f(x) - f(y) = \sum_d^{n-1} y_i + d = 6$$

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$x = 1 \ 0 \ 2 \ 0 \ 1 \ 2$$

$$x = 0 \ 2 \ 1 \ 0 \ 1 \ 2$$

$$f(x) = 24$$

$$y = 0 \ 2 \ 0 \ 1 \ 2$$

$$f(y) = 18$$

$$f(x) - f(y) = \sum_{d=1}^{n-1} y_d + d = 6$$

$$x = 1 \text{ } \boxed{0} \text{ } \boxed{2} \text{ } 0 \text{ } 1 \text{ } 2$$

$$x = 0 \text{ } \boxed{2} \text{ } \boxed{1} \text{ } 0 \text{ } 1 \text{ } 2$$

$$\sum_{d=0}^{n-1} y_i + d$$

$$\sum_{d=0}^{n-1} y_i + d$$

$$x = 1 \text{ } \boxed{0} \text{ } \boxed{2} \text{ } 0 \text{ } 1 \text{ } 2$$

$$x = 0 \text{ } \boxed{2} \text{ } \boxed{1} \text{ } 0 \text{ } 1 \text{ } 2$$

$$\sum_d^{n-1} y_i + d$$

$$\sum_d^{n-1} y_i + d$$

Fails when there is a chunk with an average 1



Weighted VT sketch

VT sketch: ternary
Why it fails

VT sketch: binary



Bias the weight!

$$f(x) = \sum_{i=1}^n i \cdot x_i$$

$$f(x) = \sum_{i=1}^n i \cdot w(x_i)$$

$$w(0) = 0$$

$$w(1) = 1$$

$$w(2) = 2 \log n$$

$$f(x) = \sum_{i=1}^n i \cdot w(x_i)$$

$$w(0) = 0$$

$$w(1) = 1$$

$$w(2) = 2 \log n$$

$$x = 1 \mathbf{0} \mathbf{2} \dots$$

$$f(x) = 1 + 6 \log n$$

$$x = \mathbf{0} \mathbf{2} 1 \dots$$

$$f(x) = 4 \log n + 3$$

As long as the chunk of avg 1 has length $< 2 \log n$

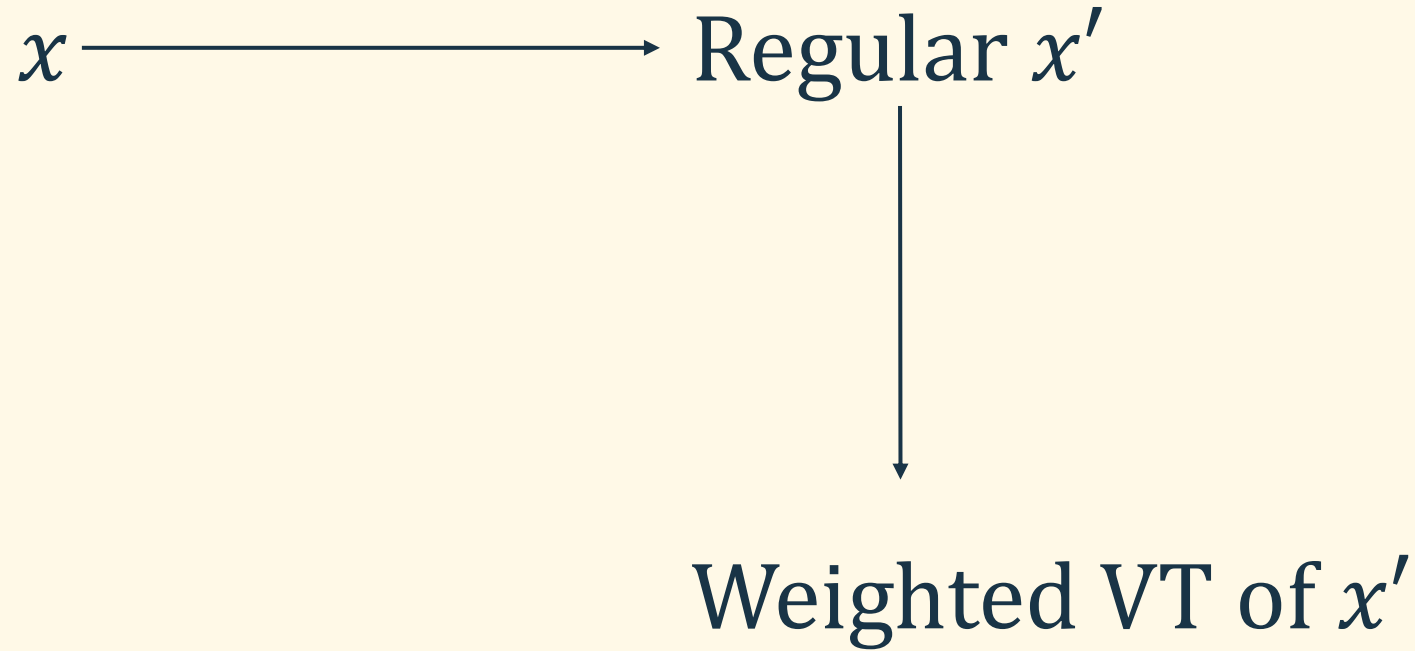
$$x = 1 \underbrace{0 \ 2 \ 0 \ 2 \ 0 \ 2 \ 0 \ 2 \ 0 \ 2}_{< 2 \log n} 2 \dots$$

Run of 0 after deleting all 2's $< \log n$

$$x = 1 \underbrace{0 \ \square \ 0 \ \square \ 0 \ \square \ 0 \ \square \ 0 \ \square}_{< \log n} 2 \dots$$

Runlength replacement [SWG17]

Weighted VT: encoding



$\log n + \log \log n$ redundancy!

More in paper

1. Extends to general q
2. Binary code correcting one del and one sub
3. Binary code correcting one del and one adjacent trans